

by  
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# Power Programming

## Creating a DPMI-Based DOS Extender Of Your Own

Life isn't easy for the authors of commercial DOS extenders these days. Their products must be capable of running on a completely standalone basis, taking full control of the 286, 386, and 486 processors in protected mode, and maintaining all the special CPU control registers, descriptor tables, and page tables. And their brainchildren must also be able to coexist harmoniously with memory managers or device drivers that implement any of the existing protocols or standards for extended memory allocation—be it top-down (INT 15h), bottom-up (VDISK), XMS (HIMEM.SYS), VCPI (QEMM-386 or 386MAX), or DPMI (Microsoft Windows 3.0). Developers of DOS extenders must also be masters of the various PC hardware architectures, with all their peculiarities and ramifications for mode switching, toggling of the A20 address line, and the like. This certainly isn't a job description I'd envy, even if it does carry its own unique rewards (the staff of Phar Lap Software has been growing, I suspect, as quickly as the Microsoft OS/2 development staff has been shrinking).

But DOS extenders are a technologically fascinating topic, and fortunately events have conspired to place a low-cost test-bed for DOS extender experimentation at our disposal. I refer, of course, to the implementation of the DOS Protected Mode Interface (DPMI) in Windows 3.0. First, let's review how "real" DOS extenders work, and then we'll implement a simple little DPMI-based DOS extender of our own.

### THE TWO FACES OF A DOS EXTENDER

A typical DOS extender that you might buy from Phar Lap or Rational Systems can be thought of as having two main sections: an initialization component and an interrupt handler component. The initialization portion of a DOS extender gains control of the machine when the DOS extender is loaded by DOS in real mode as a result of a user command or an INT

■ DOS extenders are technologically fascinating and, luckily, the Windows 3.0 implementation of the DOS Protected Mode Interface gives us a low-cost test-bed for DOS extender experimentation.

21h, function 4Bh (EXEC) function call by another program. In most cases, the DOS extender and the protected-mode application (such as *Mathematica*, *AutoCAD*, or *Lotus 1-2-3*, Release 3.0) are bound together in the same .EXE file. However, the only part of the file visible to DOS is the DOS extender, because it's the only part described by the .EXE file header and relocation table.

The initialization routine has a lot of work to do before the protected-mode application can start running. It must probe the environment and determine whether it's running on the bare hardware (aside from DOS) or in the presence of one of the several species of memory managers. It must allocate some extended memory and load the application's code and data, performing any necessary relocations and fixups. It must allocate additional conventional memory (memory below 640K) for communication with MS-DOS and the ROM BIOS. It must build the necessary global and local descriptor tables to support protected-mode addressing, which

includes mapping some descriptors onto "magic" memory areas like the PSP, the environment block, the video refresh buffer, and the ROM BIOS data area, for the convenience of the application. It must switch the CPU from real mode into protected mode. And it must take full command of the CPU's interrupt subsystem by building an interrupt descriptor table, reprogramming the 8259 interrupt controllers, and installing handlers for every type of interrupt that might occur in protected mode.

From the DOS extender's point of view, there are three types of interrupts to worry about. The first group consists of the CPU faults or exceptions, usually caused by a program error but sometimes by a hardware error (for example, a nonmaskable interrupt due to a RAM parity error) or an unexpected result on the numeric coprocessor. The second group consists of the external hardware interrupts—signals that a peripheral device has received data, is ready to accept more data, or has encountered some other condition that needs attention. The third group are the software interrupts, which occur when an application program executes an INT instruction. Control over software interrupts is a particularly crucial aspect of DOS extender initialization, because it allows the DOS extender to function transparently by intercepting the application's requests for ROM BIOS and MS-DOS services.

The exact order in which the DOS extender will perform these initialization chores depends in part on the philosophy of the implementer and in part on the

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environment in which the DOS extender will run. For example, XMS and VCPI memory managers require that the DOS extender make its memory allocation calls in real mode, while requests to a DPMI host for memory must be made in protected mode. Similarly, techniques for mode switching differ with the PC architecture (PC/AT, PS/2, and so forth) and with the presence or absence of XMS, VCPI, and DPMI. The only things we can say for sure about the initialization sequence are that a minimal global descriptor table (GDT) must be constructed in real mode and that the CPU must be in protected mode before control passes to the entry point of the actual application program.

Once the application is launched, the DOS extender's initialization section is dead code; the memory it occupies can be reclaimed and reused. The application's code now becomes the center of the action. The remaining parts of the DOS extender are activated only upon the occurrence of one of the three kinds of interrupts listed earlier. In this respect,

the DOS extender is much like a real-mode terminate-and-stay-resident (TSR) program. But the DOS extender's status after initialization is somewhat more slippery than that of your average TSR.

From MS-DOS's perspective, the DOS extender is the only program that's run-

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DOS extender can be  
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service routines.**

ning, and it's a perfectly ordinary program at that: it lives in the memory that MS-DOS allocated for it, and it makes normal MS-DOS function calls. From the protected-mode application's point of view, the DOS extender is nearly invisible (in fact, the more invisible, the better the DOS extender); the application requests

MS-DOS and ROM BIOS services by executing the usual software interrupts, and somehow the right things happen—even though the interrupts are being executed in protected mode and the application is passing virtual addresses of data and buffers lying in extended memory.

Looking at the DOS extender from the inside out (the viewpoint of its author), the fully initialized DOS extender can be thought of rather simply as a grab bag of interrupt service routines.

When a CPU fault or exception occurs, the DOS extender usually has little recourse but to terminate the protected mode application, unless the application has explicitly registered its own handler for the interrupt. If the application has gone so far awry as to cause a stack underflow or general protection (GP) fault, for example, it's obviously ailing and will only cause more trouble if it's allowed to continue. (Interpreters and debuggers are obvious exceptions.)

External hardware interrupts are treated differently. If the protected-mode application hasn't registered a handler, the interrupt is "reflected" to the original real-mode owner of the interrupt. In other words, the DOS extender fields the inter-

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; TINYDOSX.ASM Tiny DPMI-Based DOS Extender

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stdin equ 0 ; standard input handle

stdout equ 1 ; standard output handle

stderr equ 2 ; standard error handle

cr equ 0dh ; ASCII carriage return

lf equ 0ah ; ASCII line feed

DGROUP group \_DATA

\_DATA segment word public 'DATA'

mode sw dd 0 ; far pointer to DPMI mode

int0dv dd 0 ; switch entry point

int21v dd 0 ; address of previous GP fault handler

realseg dw 0 ; address of previous Int 21H handler

realseg dw 0 ; segment of real mode buffer

gpmsg db cr,lf,lf ; selector for real mode buffer

db 'TINYDOSX: general protection fault!'

db cr,lf

gpmsg len equ \$-gpmsg

abmsg db cr,lf,lf

db 'TINYDOSX: unsupported DOS function!'

db cr,lf

abmsg len equ \$-abmsg

regs label word ; real mode register structure

regDI label word ; for DPMI translation services

regDI dd 0 ; 00H DI or EDI

regSI label word ; 04H SI or ESI

regSI dd 0

regBP label word ; 08H BP or EBP

regBP dd 0

regSP label word ; 0CH (reserved)

regSP dd 0

regAX label word ; 10H AX or EAX

regAX dd 0

regBX label word ; 14H BX or EBX

regBX dd 0

regCX label word ; 18H CX or ECX

regCX dd 0

regDX label word ; 1CH DX or EAX

regDX dd 0

regEAX dd 0

cpuFlags dw 0 ; 20H CPU status flags

regES dw 0 ; 22H ES

regEBX dw 0 ; 24H EBX

regES dw 0 ; 26H ES

regEBX dw 0 ; 28H EBX

regEBX dw 0 ; 2AH EBX (CS:IP ignored by 2CH CS DPMI function 0300H)

regEBX dw 0 ; 2EH SP (SS:SP=0 to have DPMI 30H SS host supply a stack)

regEBX dw 0

protDX dw 0 ; save protected mode DX

protSI dw 0 ; save protected mode SI

protES dw 0 ; save protected mode ES

dispatch label word

dw offset TEXT:fxn00h ; fcn 00H terminate

dw offset TEXT:fxn01h ; fcn 01H char input+echo

dw offset TEXT:fxn02h ; fcn 02H char output

dw offset TEXT:fxn03h ; fcn 03H aux input

dw offset TEXT:fxn04h ; fcn 04H aux output

dw offset TEXT:fxn05h ; fcn 05H printer output

dw offset TEXT:fxn06h ; fcn 06H raw console I/O

dw offset TEXT:fxn07h ; fcn 07H raw input no echo

dw offset TEXT:fxn08h ; fcn 08H char input no echo

dw offset TEXT:abort ; fcn 09H

dw offset TEXT:abort ; fcn 0AH

dw offset TEXT:fxn0Bh ; fcn 0BH input status

dw offset TEXT:abort ; fcn 0CH

dw offset TEXT:fxn0Dh ; fcn 0DH disk reset

dw offset TEXT:fxn0Eh ; fcn 0EH select disk

dw offset TEXT:abort ; fcn 0FH

dw offset TEXT:abort ; fcn 10H

dw offset TEXT:abort ; fcn 11H

dw offset TEXT:abort ; fcn 12H

dw offset TEXT:abort ; fcn 13H

dw offset TEXT:abort ; fcn 14H

dw offset TEXT:abort ; fcn 15H

dw offset TEXT:abort ; fcn 16H

dw offset TEXT:abort ; fcn 17H

dw offset TEXT:abort ; fcn 18H

dw offset TEXT:fxn19h ; fcn 19H get current drive

dw offset TEXT:abort ; fcn 1AH

dw offset TEXT:fxn1Bh ; fcn 1BH get cur. drive data

dw offset TEXT:fxn1Ch ; fcn 1CH get drive data

dw offset TEXT:abort ; fcn 1DH

dw offset TEXT:abort ; fcn 1EH

dw offset TEXT:abort ; fcn 1FH

dw offset TEXT:abort ; fcn 20H

dw offset TEXT:abort ; fcn 21H

dw offset TEXT:abort ; fcn 22H

dw offset TEXT:abort ; fcn 23H

dw offset TEXT:abort ; fcn 24H

dw offset TEXT:abort ; fcn 25H

dw offset TEXT:abort ; fcn 26H

dw offset TEXT:abort ; fcn 27H

dw offset TEXT:abort ; fcn 28H

dw offset TEXT:abort ; fcn 29H

dw offset TEXT:fxn2Ah ; fcn 2AH get date

dw offset TEXT:fxn2Bh ; fcn 2BH set date

dw offset TEXT:fxn2Ch ; fcn 2CH get time

Figure 1: Here's the source code for a skeleton DPMI 0.9-based DOS extender that can be linked with small-model C programs.

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rupt, saves the CPU state, switches the CPU into real mode, and reissues the interrupt with an INT instruction. When the real-mode handler executes an interrupt return (IRET), the DOS extender recovers control, switches the CPU back into protected mode, restores the CPU state, and then issues its own IRET so that the application can continue.

Which brings us to the last category of interrupt that the DOS extender must dispose of: interrupts that are explicit MS-DOS or ROM BIOS function requests by the application. There are a number of different software interrupts that the DOS extender must be prepared to cope with: MS-DOS's INT 20h through INT 2Fh, the ROM BIOS video driver INT 10h, serial port driver INT 14h, keyboard driver INT 16h, printer driver INT 17h, Microsoft Mouse driver INT 33h, and so on. Each one of these interrupts provides a pathway to a host of different subfunctions, typically selected by a value in register AH. For instance, more than 80 functions (both documented and undocumented) are defined for INT 21h. As I explained in the last column, these functions can be di-

vided into four classes:

- functions that require little more than a mode switch before passing them down to MS-DOS or the ROM BIOS;
- functions that address application buffers and therefore require data movement and address translation before they can be reissued to MS-DOS or the ROM BIOS;
- functions that must be completely re-

**In any event, the DOS extender's handling of an MS-DOS or ROM BIOS service request is easy to visualize.**

placed to make them meaningful in protected mode; and

- function calls that are unique to the DOS extender itself and provide special services that have no equivalents in MS-DOS or the ROM BIOS.

There is also an implicit fifth class of functions, which I didn't mention previ-

ously: the MS-DOS functions that the DOS extender author simply chooses not to support because they're either too dangerous or not worth the hassle. The FCB file functions, the direct disk I/O functions, and some of the undocumented DOS functions that vary tremendously from one version of DOS to another are good examples.

In any event, the DOS extender's handling of an MS-DOS or ROM BIOS service request is easy to visualize. An umbrella routine is entered first; it saves the CPU flags and all the general and segment registers for later reference. The umbrella routine then decodes the function request by using the function number in AH as an index into a jump table, and then passes control to a handler that is specific to the function type. Functions that pass all parameters in registers and don't reference data in extended memory can fall through to another umbrella handler that performs the necessary mode switching and reissues the interrupt in real mode. Nearly all functions that pass addresses of buffers or pass parameters by reference must be handled on an individual basis, because there's regrettably little symmetry of structures or register usage across DOS and ROM BIOS function calls. Of course, memory management

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```

dw offset _TEXT:fxn2dh ; fxn 20h set time
dw offset _TEXT:fxn2eh ; fxn 20h set verify flag
dw offset _TEXT:fxn2fh ; fxn 20h
dw offset _TEXT:fxn30h ; fxn 30h get DOS version
dw offset _TEXT:fxn31h ; fxn 31h
dw offset _TEXT:fxn32h ; fxn 32h get/set break flag
dw offset _TEXT:fxn33h ; fxn 33h
dw offset _TEXT:fxn34h ; fxn 34h
dw offset _TEXT:fxn35h ; fxn 35h
dw offset _TEXT:fxn36h ; fxn 36h get drive info
dw offset _TEXT:fxn37h ; fxn 37h
dw offset _TEXT:fxn38h ; fxn 38h
dw offset _TEXT:fxn39h ; fxn 39h create directory
dw offset _TEXT:fxn3ah ; fxn 3Ah delete directory
dw offset _TEXT:fxn3bh ; fxn 3Bh select directory
dw offset _TEXT:fxn3ch ; fxn 3Ch create file
dw offset _TEXT:fxn3dh ; fxn 3Dh open file
dw offset _TEXT:fxn3eh ; fxn 3Eh close file
dw offset _TEXT:fxn3fh ; fxn 3Fh read file
dw offset _TEXT:fxn40h ; fxn 40h write file
dw offset _TEXT:fxn41h ; fxn 41h delete file
dw offset _TEXT:fxn42h ; fxn 42h seek
dw offset _TEXT:fxn43h ; fxn 43h get/set attributes
dw offset _TEXT:fxn44h ; fxn 44h
dw offset _TEXT:fxn45h ; fxn 45h dup handle
dw offset _TEXT:fxn46h ; fxn 46h redirect handle
dw offset _TEXT:fxn47h ; fxn 47h get cur. directory
dw offset _TEXT:fxn48h ; fxn 48h
dw offset _TEXT:fxn49h ; fxn 49h
dw offset _TEXT:fxn4ah ; fxn 4Ah
dw offset _TEXT:fxn4bh ; fxn 4Bh
dw offset _TEXT:fxn4ch ; fxn 4Ch terminate
dw offset _TEXT:fxn4dh ; fxn 4Dh
dw offset _TEXT:fxn4eh ; fxn 4Eh
dw offset _TEXT:fxn4fh ; fxn 4Fh
dw offset _TEXT:fxn50h ; fxn 50h
dw offset _TEXT:fxn51h ; fxn 51h
dw offset _TEXT:fxn52h ; fxn 52h
dw offset _TEXT:fxn53h ; fxn 53h
dw offset _TEXT:fxn54h ; fxn 54h get verify flag
dw offset _TEXT:fxn55h ; fxn 55h
dw offset _TEXT:fxn56h ; fxn 56h
dw offset _TEXT:fxn57h ; fxn 57h get/set file date
dw offset _TEXT:fxn58h ; fxn 58h
dw offset _TEXT:fxn59h ; fxn 59h
dw offset _TEXT:fxn5ah ; fxn 5Ah create temp file
dw offset _TEXT:fxn5bh ; fxn 5Bh create unique file
dw offset _TEXT:fxn5ch ; fxn 5Ch lock/unlock
dw offset _TEXT:fxn5dh ; fxn 5Dh
dw offset _TEXT:fxn5eh ; fxn 5Eh
dw offset _TEXT:fxn5fh ; fxn 5Fh
dw offset _TEXT:fxn60h ; fxn 60h
dw offset _TEXT:fxn61h ; fxn 61h
dw offset _TEXT:fxn62h ; fxn 62h
dw offset _TEXT:fxn63h ; fxn 63h
dw offset _TEXT:fxn64h ; fxn 64h
dw offset _TEXT:fxn65h ; fxn 65h
dw offset _TEXT:fxn66h ; fxn 66h
dw offset _TEXT:fxn67h ; fxn 67h
dw offset _TEXT:fxn68h ; fxn 68h commit file
dw offset _TEXT:fxn69h ; fxn 69h

dw offset _TEXT:abort ; fxn 6Ah
dw offset _TEXT:fxn6bh ; fxn 6Bh
dw offset _TEXT:fxn6ch ; fxn 6Ch
dw offset _TEXT:fxn6dh ; fxn 6Dh
dw offset _TEXT:fxn6eh ; fxn 6Eh
dw offset _TEXT:fxn6fh ; fxn 6Fh

DATA ends
_TEXT segment byte public 'CODE'
    assume cs:_TEXT,ds:DGROUP
;
; Initialization routine for the Tiny DOS Extender. First we test for
; the presence of a DPMI host, get the address of the mode switch entry
; point, and request the switch to protected mode. Then we install
; a handler for GP faults to circumvent the Win 3 brain-damaged dialog
; box, and allocate some memory below 1 MB to use as a buffer for
; communication with DOS. Finally we install our own int 21h handler
; so we can service DOS calls from the protected mode application.
;
    public initdosx
initdosx proc near
    mov ax,1697h ; get address of DPMI
    int 2fh ; mode switch entry point
    or ax,ax ; bail out if no DPMI
    jnz init9
    mov word ptr modesw,di ; save far pointer to
    mov word ptr modesw+2,es ; DPMI entry point

    mov bx,si ; allocate DPMI private data
    mov ah,40h ; area below 1 MB boundary
    int 21h
    jc init9 ; jump, allocation failed

    mov es,ax ; pass segment of data area
    mov ax,0 ; bit 0=0 indicates 16-bit app
    call modesw ; switch to protected mode

    mov ax,0202h ; get address of previous
    mov bl,dh ; owner of GP fault vector
    int 31h
    mov word ptr int8dv,dx ; save as far pointer
    mov word ptr int8dv+2,cx

    mov ax,0203h ; install our GP fault handler
    mov bl,dh
    mov cx,0e ; CX:DX = handler address
    mov dx,offset _TEXT:upfpiar
    int 31h
    jc init9 ; jump, couldn't install

    mov ax,0100h ; allocate 64 KB buffer in
    mov bx,1000h ; conventional memory for
    int 31h ; communication with DOS
    jc init9 ; jump, allocation failed
    mov realseg,ax ; save segment of block
    mov realseal,dx ; save selector for block

    mov ax,0204h ; get address of previous

```

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function calls must be trapped and serviced entirely within the DOS extender, and this can be a fair amount of work in itself (particularly if the DOS extender supports virtual memory).

One more function that bears mentioning is INT 21h, function 4Ch, which an application calls to terminate itself. When the DOS extender sees this function request, it must deallocate its (and the application's) extended memory, release the interrupt subsystem, and remove any other little tendrils it may have inserted elsewhere. This is tricky business, because the DOS extender—which has been playing the role of a little protected-mode operating system that uses DOS as a file system slave—must quietly and gracefully put everything back *exactly* as it was before the DOS extender was loaded, leaving the system completely stable. Finally, the DOS extender must switch the CPU into real mode for the last time and itself call INT 21h, function 4Ch, so that DOS knows to free all conventional memory, file handles, and other resources that were assigned to the DOS extender.

## A SIMPLE DPMI-BASED DOS EXTENDER

I always find source code much more enlightening than windy explanations, so I've illustrated the foregoing discussion of DOS extenders with TINYDOSX.ASM (Figure 1), a skeleton DOS extender that you can use in your own programs. TINYDOSX relies only on the existence

**Memory management function calls must be trapped and serviced entirely within the DOS extender, which can be a fair amount of work.**

of a DPMI 0.9 host, such as the one found in *Windows 3.0*; it can be linked into any small-model C program and will cause that program to execute in protected mode—provided you don't call any run-time library functions that include self-

modifying code, use segment registers for scratch storage, perform segment register arithmetic, or in general execute an instruction that will result in a general protection fault. Although *Windows 3.0* includes a DOS extender of its own, TINYDOSX doesn't use it and will run just as well in other DPMI, Version 0.9, environments (at least theoretically; no other such environments exist for testing at this time).

The initialization portion of TINYDOSX, embodied in the routine INITDOSX, is straightforward. To keep things simple, we allow the C application to get control first in real mode, and require it to explicitly call the DOS extender, rather than the other way around. INITDOSX first calls INT 2Fh, function 1687h to find out whether a DPMI host is present, and if so, the address of the mode switch entry point. If a DPMI host is not found, INITDOSX bails out with an error message; otherwise, it proceeds to allocate the private data area required by the DPMI host and then requests the switch into protected mode.

Once it's running in protected-mode, INITDOSX installs protected mode handlers for MS-DOS INT 21h and for general protection faults (so that *Windows* won't

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```

mov     bl,21h           ; owner of int 21h vector
int     21h
mov     word ptr int21v,dx ; save as far pointer
mov     word ptr int21v+2,cx

mov     ax,0285h         ; install our int 21h handler
mov     bl,21h
mov     cx,cx
mov     dx,offset _TEXT:doscalle
int     21h
xor     ax,ax
ret

init9:  mov     ax,-1
ret

initdosx endp

; Interrupt service routine for GP faults. Entered by a far call
; from DPMI host with CS:IP, flags, CPU error code on stack.
; We force transfer to our error message routine by changing
; return address in the stack frame.
;
; gpfiar proc     far
;
;     push     bp
;     mov     bp,sp
;     mov     word ptr [bp+8],offset _TEXT:gpfierr
;     pop     bp
;     ret
;
; gpfiar endp
;
; This routine gains control after the GPFIAR returns to the DPMI host.
; It simply displays an error message and terminates cleanly, subverting
; Win 3's "Application has violated system integrity" dialog box.
;
; gpfierr proc     near
;
;     mov     dx,offset DGROUP:gpfierrmsg ; display GP fault message
;     mov     cx,gpfierrmsg_len
;     mov     bx,stdout
;     mov     ah,40h
;     int     21h
;     mov     ax,4c01h
;     int     21h
;
; gpfierr endp
;
; The DOSCALLE routine is the runtime portion of the Tiny DOS Extender.
; It traps int 21h requests in protected mode and performs any necessary
; mode switching, data movement, and address translation on a
; function-by-function basis. Anything DOSCALLE doesn't want to handle,
; it either fails by setting the carry flag and returning, or
; it aborts the current program. In particular, all FCB-related
; functions are aborted. When a termination function is detected,
; the interrupt handlers are unhooked and the function call is
; passed down to the DPMI host so that all other protected mode
; resources will be deallocated.

```

```

doscalle proc     far
;
;     push     bx
;     mov     bl,ah
;     xcr     bh,bh
;     cmp     bx,6fh
;     ja      abort
;     add     bx,bx
;     jmp     [dispatch+bx]
;
; abort:
;     mov     dx,offset DGROUP:abrtmsg
;     mov     cx,abrtmsg_len
;     mov     bx,stdout
;     mov     ah,40h
;     int     21h
;     mov     ax,4c01h
;     int     21h
;
; error:
;     push     bp
;     mov     bp,sp
;     or     word ptr [bp+6],1
;     pop     bp
;     mov     ax,1
;     iret
;
; common handling for entirely
; register-based functions
; function 01h: char input+echo
; function 02h: char output
; function 03h: aux input
; function 04h: aux output
; function 05h: printer output
; function 06h: raw console I/O
; function 07h: raw input no echo
; function 08h: char input no echo
; function 09h: input status
; function 0Ah: disk reset
; function 0Bh: select disk
; function 0Ch: get current drive
; function 0Dh: get cur. drive data
; function 0Eh: get drive data
; function 0Fh: get date
; function 10h: set date
; function 11h: get time
; function 12h: set time
; function 13h: set verify flag
; function 14h: get DOS version
; function 15h: get/set break flag
; function 16h: get drive info
; function 17h: close file
; function 18h: seek
; function 19h: dup handle
; function 1Ah: redirect handle
; function 1Bh: get verify flag
; function 1Ch: get/set filedate
; function 1Dh: lock/unlock
; function 1Eh: commit file
;
;     pop     bx
;     call    saveregs
;     unld    general registers

```



```

call    realdos      ; transfer to DOS
call    loadregs     ; load general registers & flags
iret                     ; return to application

; common handling for functions
; passing ASCIIIZ addr in DS:DX
; function 39H: create directory
; function 3AH: delete directory
; function 3BH: select directory
; function 3CH: create file
; function 3DH: open
; function 41H: delete file
; function 43H: get/set attributes
; function 5AH: create temp file
; function 5BH: create unique file
; restore BX
; unload general registers
; ES:DI = virtual address of
; real mode buffer
; DS:SI = virtual address of
; protected mode buffer
; copy ASCIIIZ string to
; real mode buffer
; reached null yet?
; no, copy another character
; set address of real mode buffer
; into register data structure
; transfer to MS-DOS
; load general registers & flags
; restore protected mode DX, ES
; return to application

; function 3FH: read file
; restore BX
; unload general registers
; set address of real mode buffer
; into register data structure
; transfer to MS-DOS
; CX = actual length of data
; ES:DI = virtual address of
; protected mode buffer
; DS:SI = virtual address of
; real mode buffer
; copy data from real mode
; buffer to protected mode buffer
; restore DS = our DGROUP
; load general registers
; restore protected mode DX, ES
; return to application

; function 40H: write file
; restore BX
; unload general registers
; ES:DI = virtual address of
; real mode buffer
; DS:SI = virtual address of
; protected mode buffer
; copy data to real mode buffer
; set address of real mode buffer
; into register data structure
; transfer to MS-DOS
; load general registers & flags
; restore protected mode DX, ES
; return to application

; function 47H: get directory
; restore BX
; unload general registers
; set address of real mode buffer
; into register data structure
; transfer to MS-DOS
; ES:DI = virtual address of
; protected mode buffer
; DS:SI = virtual address of
; real mode buffer
; copy ASCIIIZ string from real
; mode buffer to prot mode buffer
; found null character yet?
; no, copy another character
; restore DS = our DGROUP
; load general registers
; restore protected mode SI, ES
; return to application

; function 00H: terminate
; function 4CH: terminate
; save return code
; restore old GP fault handler
; restore old Int 21H handler
; release real mode buffer
; chain to DPMI Int 21H handler
; for cleanup and termination
; general fallback point
; (useful during debugging)
; restore register BX
; chain to prev Int 21H owner

doscall endp

; Save general registers into real mode data structure for a call to
; real mode routine via DPMI translation services. Note that segment
; registers are NOT unloaded into structure because they are not valid
; for real mode anyway.
;
; saveregs proc    near
;
;     mov     regAX,ax
;     mov     regBX,bx
;     mov     regCX,cx
;     mov     regDX,dx
;     mov     regSI,si
;     mov     regDI,di
;     mov     regBP,bp
;     mov     protDX,dx
;     mov     protSI,si
;     mov     protES,es
;     ret
;
; saveregs endp
;
; Load general registers from real mode data structure. Note that
; segment registers are NOT loaded because their real mode values
; would cause a GP fault in protected mode.
;
; loadregs proc    near
;
;     mov     bp,sp
;     push    cpuFLAGS
;     pop     [bp+6]
;     mov     ax,regAX
;     mov     bx,regBX
;     mov     cx,regCX
;     mov     dx,regDX
;     mov     si,regSI
;     mov     di,regDI
;     mov     bp,regBP
;     ret
;
; loadregs endp
;
; Call the DPMI translation function 0300H to simulate a real mode
; software interrupt 21H, transferring control to MS-DOS, passing
; the values stored into the real mode register structure 'regs'.
;
; realdos proc    near
;
;     push    es
;     mov     ax,0300h
;     mov     bl,21h
;     mov     bh,0
;     mov     cx,8
;     push    ds
;     mov     es,di
;     mov     di,offset DGROUP:regs
;     int     31h
;     pop     es
;     ret
;
; realdos endp
;
; _TEXT ends
;
; end

```

blow our little programming blunders out of the water with its totally uninformative "Application has violated system integrity" dialog box). INITDOSX also allocates a 64K area of conventional memory that the INT 21h handler can later use to pass data back and forth to MS-DOS. Finally, INITDOSX returns control to the C application, which continues its execution in protected mode.

When the C application requests an MS-DOS service, the protected-mode INT 21h handler, named DOSCALL, receives control. DOSCALL saves the flags and

registers, then branches through the table DISPATCH to the appropriate subroutine. As you'll notice, I've stubbed out most of the less-common DOS functions to either return an error or abort the application. The functions that TINYDOSX supports, however, are relayed to DOS using the DPMI translation function "Simulate Real Mode Interrupt." Use of this translation function, rather than the speedier DPMI "raw mode switch" function, eliminates all sorts of messy problems that are best left to the imagination and experiments of adventurous readers.

Naturally, DOSCALL monitors for the fateful INT 21h, function 4Ch, and cleans up after itself accordingly.

Assuming DPMI 0.9 as our platform and linking TINYDOSX directly into the protected-mode application allows us to take some shortcuts that would never suffice in a commercial-grade DOS extender. First, we don't have to build our own loader for the protected-mode application; since TINYDOSX is linked into the application, DOS loads TINYDOSX as the application as a single unit. Second, we don't have to allocate any memory

## Power Programming

### TINYDOSX.C

```

/*
TESTDOSX.C --- illustrates use of the DPMI-based
DOS Extender TINYDOSX to display a message in protected mode

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PC Magazine * Ray Duncan

Build with Microsoft C 6.0 SMALL MODEL as follows:

    MASM TINYDOSX;
    CL TESTDOSX.C TINYDOSX

Execute under Windows 3.0 in the DOS Box only!
*/

#include <stdio.h>
unsigned extern pascal InitDosX(void);
main()
{
    unsigned saveCS, saveDS;
    _asm mov saveCS,cs          ; store real mode CS
    _asm mov saveDS,ds          ; and DS for display
    printf("\nHello, real mode world! \tcs=%04xh DS=%04xh",
        saveCS, saveDS);

    if(InitDosX())              // attempt mode switch
    {
        puts("\nDPMI initialization failed.");
        exit(1);
    }

    _asm mov saveCS,cs          ; store protected mode CS
    _asm mov saveDS,ds          ; and DS for display
    printf("\nHello, protected mode world! \tcs=%04xh DS=%04xh\n",
        saveCS, saveDS);

    _asm mov ah,4ch              ; exit directly to DOS to avoid
    _asm int 21h                 ; GP fault in RTL cleanup code
}

```



**Figure 2:** This is the source code for a simple protected-mode C application that can be linked with TINYDOSX.ASM, listed in Figure 1.

for the application or build any descriptors; we get these services “free” when DPMI creates code and data selectors during the initial switch to protected mode. Third, we need to support only those DOS services in our DOS extender that we know our application is actually going to use; we don’t have the obligation to translate every known (and unknown) DOS function for protected mode the way a “real” DOS extender does.

Suppose we wanted to turn TINYDOSX into a not-so-tiny, more robust DOS extender—where would we start? We’d have to enlarge the support for INT 21h functions to include (at minimum) all the documented MS-DOS services. We’d need to support the immense battery of ROM BIOS services (most of which, luckily, are register-based anyway) and probably, in addition, the Microsoft Mouse INT 33h and the NETBIOS interfaces. We’d want to add more sophisticated facilities for installation of interrupt handlers by the application. And last but not least, we’d be obligated to support most of the different C memory models, a somewhat arduous chore. My first impression of the way to do this would be to incorporate a loader for “segmented .EXE” (also called

“new .EXE”) files into our DOS extender, build the application with the Microsoft Segmented Linker, and bind our DOS extender into the .EXE file as the “real-mode stub.”

#### TRYING OUT TINYDOSX

TESTDOSX.C, a simple protected-mode C application that can be linked with TINYDOSX.ASM, is shown in Figure 2. To create the executable version of TESTDOSX, enter the following commands:

```

MASM /Mx TINYDOSX;
CL TESTDOSX.C TINYDOSX

```

Be sure that you are using Microsoft 6.0 and compiling for the small model. The resulting application, TESTDOSX.EXE, must be run in one of Windows 3.0’s DOS boxes so that it has access to DPMI services.

#### THE IN-BOX

Please send your questions, comments, and suggestions to me at any of the following e-mail addresses:  
 PC MagNet: 72241,52  
 MCI Mail: rduncan  
 BIX: rduncan

## Computers, FCC Class A, Class B, and You — or When is it better to get a B than an A?

You need to know the difference between computers that meet the FCC class B radio frequency emissions standards and those that meet only the Class A standards.

Computers emit radio signals in their operation. Because these signals may cause interference to radio and television reception, the marketing and the use of computers is regulated by the Federal Communications Commission. Under federal rules, computer users are responsible for remedying interference, including interference in neighboring homes.

Computers certified by the FCC as meeting the Class B standard are less likely to cause interference to radio and TV reception than those that have been verified by the manufacturer or importer to the Class A standards. Only Class B certified computers may be advertised, sold, or leased for use in residences. A similar regulatory program applies in Canada.

Buyers seeking computers for use in homes (including offices at home) should shop for computers and peripherals which have been Class B certified. These devices carry a label with an FCC ID number. Both new and used Class A verified devices may be sold only for use in commercial and industrial locations. Signals from computers are more likely to be masked by electrical noise from other equipment in such an environment. These areas are also likely to have fewer radios and TVs. Accordingly, equipment marketed only for use in these locations may meet the less rigorous Class A standard. Class B certified equipment may be marketed for use in residences as well as commercial and industrial locations.

As you shop for a computer for use in your home, look for the FCC classification in the specifications or ask your vendor to recommend only machines that have been certified to the Class B limits. TV viewers and radio listeners in your home and in neighboring homes will be glad you did.